

FORM PTO-1390		U.S. DEPARTMENT OF COMMERCE PATENT AND TRADEMARK OFFICE TRANSMITTAL LETTER TO THE UNITED STATES DESIGNATED/ELECTED OFFICE (DO/EO/US) CONCERNING A FILING UNDER 35 U.S.C. 371	ATTORNEY'S DOCKET NUMBER: 53863-64312
INTERNATIONAL APPLICATION NO.: PCT/SE99/02244		INTERNATIONAL FILING DATE: 01 December 1999	PRIORITY DATE CLAIMED: 01 December 1998

TITLE OF INVENTION: ACCESS CONTROL MECHANISM FOR PACKET SWITCHED COMMUNICATION NETWORKS

APPLICANT(S) FOR DO/EO/US: Christer LANDBERG and Lars OLGUS

Applicant herewith submits to the United States Designated/Elected Office (DO/EO/US) the following items and other information:

1. ☒ This is a **FIRST** submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a **SECOND** or **SUBSEQUENT** submission of items concerning a filing under 35 U.S.C. 371.
3. ☒ This express request to begin national examination procedures (35 U.S.C. 371(f)) at any time rather than delay examination until the expiration of the applicable time limit set in 35 U.S.C. 371(b) and PCT Articles 22 and 39(1).
4. ☒ A proper Demand for International Preliminary Examination was made by the 19th month from the earliest claimed priority date.
5. ☒ A copy of the International Application as filed (35 U.S.C. 371(c)(2))
 - a. ☒ is transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ has been transmitted by the International Bureau. (see attached copy of PCT/IB/308)
 - c. ☐ is not required, as the application was filed in the United States Receiving Office (RO/US).
6. ☐ A translation of the International Application into English (35 U.S.C. 371(c)(2)).
7. ☐ Amendments to the claims of the International Application under PCT Article 19 (35 U.S.C. 371(c)(3)).
 - a. ☐ are transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ have been transmitted by the International Bureau.
 - c. ☐ have not been made; however, the time limit for making such amendments has NOT expired.
 - d. ☐ have not been made and will not be made.
8. ☐ A translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 371(c)(3)).
9. ☐ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)).
10. ☐ A translation of the annexes of the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Item 11. to 16. below concern document(s) or information included:

11. ☒ An Information Disclosure Statement under 37 CFR 1.97 and 1.98.
12. ☐ An assignment document for recording. A separate cover sheet in compliance with 37 CFR 3.28 and 3.31 is included.
13. ☒ A **FIRST** preliminary amendment.
14. ☐ A **SECOND** or **SUBSEQUENT** preliminary amendment.
15. ☐ A substitute specification.
16. ☒ A change of power of attorney and/or address letter.
17. ☐ Other items or information:

International Preliminary Examination Report (PCT/IPEA/409)
International Search Report (PCT/ISA/210)
Abstract on a separate sheet
Application Data Sheet

U.S. APPLICATION NO. 09/057288

INTERNATIONAL APPLICATION NO.
PCT/SE99/02244ATTORNEY'S DOCKET NO.
53863-6431217. ☒ The following fees are submitted:**BASIC NATIONAL FEE (37 CFR 1.492(a)(1)-(5)):**

Neither international preliminary examination fee (37 CFR 1.482) nor international search fee (37 CFR 1.445(a)(2)) paid to USPTO and International Search Report not prepared by the EPO or JPO

\$ 1,000.00

International preliminary examination fee (37 CFR 1.482) not paid to USPTO but International Search Report prepared by the EPO or JPO

\$ 860.00

International preliminary examination fee (37 CFR 1.482) not paid to USPTO but international search fee (37 CFR 1.445(a)(2)) paid to USPTO

\$ 710.00

International preliminary examination fee (37 CFR 1.482) paid to USPTO but all claims did not satisfy provisions of PCT Article 33(1)-(4)

\$ 690.00

International preliminary examination fee (37 CFR 1.482) paid to USPTO and all claims satisfied provisions of PCT Article 33(1)-(4)

\$ 100.00

ENTER APPROPRIATE BASIC FEE AMOUNT =

\$ 1,000.00

Surcharge of \$130.00 for furnishing the oath or declaration later than 30 months from the earliest claimed priority date (37 CFR 1.492(e)).

\$ 130.00

CLAIMS	NUMBER FILED	NUMBER EXTRA	RATE	
Total claims	20 - 20 =	0	X \$18.00	\$
Independent claims	8 - 3 =	5	X \$80.00	\$ 400.00

MULTIPLE DEPENDENT CLAIMS(S) (if applicable)

+ \$270.00

TOTAL OF ABOVE CALCULATIONS =

\$ 1,530.00

Reduction of 1/2 for filing by small entity, if applicable. Applicant claims Small Entity Status under 37 CFR 1.27. +

SUBTOTAL =

\$ 1,530.00

Processing fee of \$130 for furnishing the English translation later than months from the earliest claimed priority date (37 CFR 1.49(f)).

TOTAL NATIONAL FEE =

\$ 1,530.00

Fee for recording the enclosed assignment (37 CFR 1.21(h)). The assignment must be accompanied by an appropriate cover sheet (37 CFR 3.28, 3.31). \$40.00 per property +

TOTAL FEES ENCLOSED =

\$ 1,530.00

Amount to be
refunded:

charged:

a. ☒ A check in the amount of \$ 1,530.00 to cover the above fees is enclosed.b. ☐ Please charge my Deposit Account No. **25-0120** in the amount of \$ to cover the above fees. A duplicate copy of this sheet is enclosed.c. ☒ The Commissioner is hereby authorized to charge any additional fees which may be required by 37 CFR 1.16 and 1.17, or credit any overpayment to Deposit Account No. **25-0120**. A duplicate copy of this sheet is enclosed.

SEND ALL CORRESPONDENCE TO:

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June 1, 2001

By

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PATENTS

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

In re application of

Christer LANDBERG et al.

Box PCT

Serial No. (unknown)

Application Branch

Filed herewith

ACCESS CONTROL MECHANISM
FOR PACKET SWITCHED
COMMUNICATION NETWORKS

PRELIMINARY AMENDMENT

Commissioner for Patents

Washington, D.C. 20231

Sir:

Prior to the first Official Action and calculation of the filing fee, please amend the above-identified application as follows:

IN THE CLAIMS:

Amend claim 3 as follows:

--3. (Amended) A method as claimed in claim 1, characterized by allocating said time slot to a further network node (20) carrying a non-delay sensitive service on the basis of stored information concerning non-delay sensitive traffic awaiting transmission if the scheduler (30, 40) indicates that no scheduling interval has elapsed.--

Amend claim 4 as follows:

--4. (Amended) A method as claimed in claim 1, characterised by determining from said scheduler (30, 40) if the scheduling interval for more than one traffic service has

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elapsed, and allocating consecutive time slots to the scheduled traffic services with alternating priority.--

Amend claim 9 as follows:.

--9. (Amended) A network as claimed in claim 7, characterised in that the cells of the scheduling means are accessed cyclically.--

Amend claim 10 as follows:

--10. (Amended) A network as claimed in claim 7, characterised in that the number of cells included in said scheduling means (30) is at least equal to the multiple of the scheduling intervals in terms of time slots of the scheduled traffic.--

Amend claim 11 as follows:

--11. (Amended) A network as claimed in claim 7, characterised in that said scheduling means (40) comprises several cyclical schedulers (41), wherein each cyclical scheduler (41) is programmed with scheduling markers relating to traffic services having the same scheduling interval in terms of time slots, the cells of each scheduler corresponding to the same time slot being accessible in turn.--

Amend claim 13 as follows:

--13. (Amended) a network as claimed in claim 6, characterised in that said central node (10) comprises queue recording means (11) for storing the packet queue size of non-delay sensitive traffic awaiting transmission at said network nodes (20).--

R E M A R K S

The above changes in the claims merely place this national stage application in the same condition as it was during Chapter II of the international stage, with the multiple dependencies being removed.

Attached hereto is a marked-up version of the changes made to the claims by the current amendment. The attached page is captioned "VERSION WITH MARKINGS TO SHOW CHANGES MADE."

Respectfully submitted,

YOUNG & THOMPSON

By

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June 1, 2001

VERSION WITH MARKINGS TO SHOW CHANGES MADE

Amend claim 3 as follows:

--3. (Amended) A method as claimed in claim 1 ~~or 2~~, characterized by allocating said time slot to a further network node (20) carrying a non-delay sensitive service on the basis of stored information concerning non-delay sensitive traffic awaiting transmission if the scheduler (30, 40) indicates that no scheduling interval has elapsed.--

Amend claim 4 as follows:

--4. (Amended) A method as claimed in ~~any previous~~ claim 1, characterised by determining from said scheduler (30, 40) if the scheduling interval for more than one traffic service has elapsed, and allocating consecutive time slots to the scheduled traffic services with alternating priority.--

Amend claim 9 as follows:

--9. (Amended) A network as claimed in claim 7 ~~or 8~~, characterised in that the cells of the scheduling means are accessed cyclically.--

Amend claim 10 as follows:

--10. (Amended) A network as claimed in ~~any one of~~ claims ~~7 to 9~~, characterised in that the number of cells included in said scheduling means (30) is at least equal to the multiple of the scheduling intervals in terms of time slots of the scheduled traffic.--

Amend claim 11 as follows:

--11. (Amended) A network as claimed in ~~any one of~~ claims 7 ~~to 9~~, characterised in that said scheduling means (40) comprises several cyclical schedulers (41), wherein each cyclical scheduler (41) is programmed with scheduling markers relating to traffic services having the same scheduling interval in terms of time slots, the cells of each scheduler corresponding to the same time slot being accessible in turn.-

Amend claim 13 as follows:

--13. (Amended) a network as claimed in ~~any one of~~ claims 6 ~~to 12~~, characterised in that said central node (10) comprises queue recording means (11) for storing the packet queue size of non-delay sensitive traffic awaiting transmission at said network nodes (20).--

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In a point-to-multipoint link traffic is transmitted in packets across the shared medium. Access to the medium is controlled by permits sent by a central node to multipoint nodes in response to a queue record of the number of packets awaiting transmission at the multipoint nodes. When circuit emulated traffic, or traffic with tight delay and delay variation requirements, is to be sent across the link, permits are sent by the central node unsolicited at regular intervals in accordance with information stored in a scheduler. The scheduler may be a single or several circular memories comprising cells representing time slots on the link. The cells contain scheduling information corresponding to one or more multipoint nodes. The central node consults the scheduler to determine if a permit is scheduled. If no permit is scheduled the time slot may be allocated to an unscheduled service on the basis of a queue record.

2/PRTS

An access control mechanism for packet switched communication networks

1. Field of invention

5 The invention relates to media access control (MAC) over a shared multiplexed link using packet transmission. It is particularly directed to the management of the transfer of information from services which stipulate certain delays and delay variations, such as circuit emulated traffic, in a packet switched network, and specifically between a central node and one of several multipoint nodes over a shared medium.

2. Background art

Media access control (MAC) in a packet switched network is a mechanism for determining which of a plurality of nodes is allowed to send information over a shared medium. In a point-to-multipoint configuration it is common that the head end or central node in a shared link will control the utilisation of the link by multipoint nodes. This is achieved by polling the various multipoint nodes. The polling order and frequency is generally controlled according to traffic demand. In a time division multiplexed TDM link, time slots are allocated by the head end of a link in response to requests sent by the multipoint nodes when traffic is awaiting transmission at these node. The slots are allocated using permits that designate which time slot in the future traffic stream may be utilised for which service traffic. When more traffic is waiting than there are available time slots, the central node will issue queuing and prioritising requests to determine which service should be allocated channel capacity first, and which can wait.

Packet switched networks are increasingly utilised to transport a variety of traffic, some of which is traditionally circuit switched. An example of this is the transmission of voice over an internet protocol (IP) based network. While speech information may be packetised, certain transmission delays and delay variations must be respected if the speech service is to be acceptable to the end user. This is a general requirement of all circuit emulated traffic. Accordingly, with the above-described request based protocols, a multipoint node carrying circuit emulated traffic requires permits to be allocated without delay and with high priority. When more than one delay sensitive service utilises the shared medium the access control becomes more complex. This situation is made still more difficult when the delay sensitive services use different frame lengths and possibly also different bit rates so that intervals between the required high priority permits are always changing. The result is often excessive delays in one or more of the delay sensitive services leading to the end user experiencing an unacceptable quality of service.

SUMMARY OF INVENTION

It is therefore an object of the invention to provide a method and arrangement enabling the efficient handling of hybrid traffic over a shared multiplexed link.

It is a further object of the present invention to provide a method and arrangement enabling the handling of several delay sensitive services over a shared medium in a packet switched network.

These and further objects are achieved in a method and arrangement for controlling access to a central node over shared medium by one of several network nodes carrying both delay sensitive and non-delay sensitive traffic, by

providing a scheduler that defines a predetermined scheduling interval for the delay sensitive traffic services. The scheduling interval corresponding essentially to the time in terms of time slots for assembling a data packet of the delay sensitive traffic service. Prior to allocating a time slot, the central node consults the scheduler to determine if a scheduling interval has elapsed. If an interval has elapsed, a time slot is allocating to the network node carrying the corresponding delay sensitive traffic service. If no scheduling interval has elapsed for the time slot consulted, the central node allocates a time slot to a non-delay sensitive service on the basis of stored information concerning the amount of traffic awaiting transmission at the various network nodes.

The scheduler is a cyclical storage means that is divided into cells representing a time slot on the shared medium and adapted to contain scheduling information relating to the traffic service.

The scheduler may comprise scheduling information for several services in a single cyclical storage means. Alternatively the scheduler contains several separate, cyclical schedulers arranged in a queue, and each being dedicated to a single traffic service.

With this method and arrangement, the transmission efficiency across the shared link can be substantially improved because no exchange of queue information is necessary for the scheduled services. Furthermore, the delays and delay variations experienced by circuit emulated and other delay sensitive traffic can be kept at a minimum. By using multiple cyclical storage means, one dedicated to each service, both the implementation of the scheduler and its expansion to accommodate additional delay sensitive services is rendered exceptionally simple.

BRIEF DESCRIPTION OF THE DRAWINGS

Further objects and advantages of the present invention will become apparent from the following description of the preferred embodiments that are given by way of example with reference to the accompanying drawings, in which:

Fig. 1 schematically depicts a point to multipoint link of a multiplexed packet switched network,

Fig. 2 shows a first embodiment of a scheduler according to the invention, and

Fig. 3 shows a further embodiment of a scheduler according to the invention.

DETAILED DESCRIPTION OF THE DRAWINGS

A typical multipoint-to-point link in a TDMA packet switched communication network is illustrated in Fig. 1. Several multipoint nodes 20 communicate via a single, shared bi-directional channel with a central node 10. Only three nodes are shown in the figure, however it is to be understood that further nodes may be connected above and below the illustrated nodes. The channel 100 may be a fixed electrical or electro-optical link or an air interface.

As in many conventional media access control (MAC) protocols, the central node 10 determines which time slots on the shared channel 100 will be allocated to which services conveyed by the multipoint nodes 20. This is accomplished by the central node 10 issuing permits for a time slot in the future information stream.

For delay insensitive packet switched services, such as data transmission services, the central node 10 holds a record of the number of packets of information awaiting transmission at each of the multipoint nodes 20. This is represented in the figure by a queue record memory 11. The queue records are updated by means of permit requests, or equivalent queue size reports, sent by the multipoint nodes 20 to the central node 10. When the record of a particular multipoint node 20 indicates that packets are awaiting transmission, the central node 10 issues a corresponding permit defining which time slot the multipoint node 20 in question may utilise for transmitting the packet. When more packets are awaiting transmission than can be dealt with immediately, the central node 10 determines which service has priority. This prioritising information is likewise obtained from the multipoint nodes 20 on invitation from the central node 10.

In the present example, however, it is assumed that the traffic A, B, C conveyed by the three illustrated multipoint nodes 20 is delay sensitive, i.e. has strict requirements as to allowable delays and delay variations. This may include, but is not limited to, circuit emulated traffic, voice or interactive video traffic. Furthermore, each service A, B, C has different frame lengths.

In accordance with the present invention, these services are not handled in the same way as delay insensitive services. While the central node 10 issues permits to the services A, B, C, these are not issued in response to information obtained from the multipoint nodes 20, but are instead scheduled by the central node 10 at regular intervals. In other words, no exchange of information takes place between the nodes. The issuing of permits is unsolicited. This is performed in accordance with the present invention with the aid of a scheduler.

Fig. 2 shows a scheduler 30 according to a first embodiment of the invention. The scheduler 30 can be viewed as a circular memory divided into cells 31, each of which represents a time slot on the shared channel. It will be understood, however, that the scheduler 30 may be implemented using one or more modulo counters, memories or registers with pointers, or similar means which allow information to be obtained at predetermined regular intervals. The scheduler 30 is either incorporated in the central node 10 or is arranged such that a controller of the central node 10 (not shown) that is adapted to handle the allocation of time slots across the shared link can easily access the scheduler 30.

Prior to operation, or upon the introduction of a new delay sensitive service, the scheduler 30 is programmed. This may be done by the central node 10 using its controller, or may be performed using external means. The scheduler 30 in Fig. 2 is already programmed with markers in some of the cells 31 that serve to indicate which service is to be allocated a permit for which time slot in the future traffic stream. In the present example these markers are represented by the letters A, B and C for the three delay sensitive services of the same names, respectively. The scheduler 30 does not define the slot as an absolute time, but rather the slot interval. The different services, A, B and C, have different frame formats and lengths determined by the transmission bit rate and denoted in the figure by a, b and c, respectively. In the figure, the frame length a of the service A is given as the equivalent of 7 time slots on the shared channel; the frame length b of service B is 8 time slots, and the frame length c of the service C is 11 time slots. In order that the permit intervals for each of these services can be respected, the total length of the scheduler 30 is a multiple of all the frame lengths in terms of time slots. While the scheduler 30 illustrated in Fig. 2 is somewhat truncated, it is assumed to have a length equal to the smallest multiple of a, b and c, that is 616 cells 31. It will be understood

that the permit intervals between different services will also take account of possible path delay differences between the various multipoint nodes 20 and the central node 10.

5 In operation, the central node 10, or its controller, consults the cells 31 of scheduler 30 in succession to determine whether a permit is scheduled for one of the services A, B or C for each time slot. If the cell 31 consulted is marked, a permit will be issued to the corresponding multipoint node 20 carrying the selected service. If no permit is scheduled for a particular time slot, the central
10 node 10 may issue a permit to one of the other delay insensitive services which require access to the shared medium on the basis of queue size information stored at the node 10, as described above. Once all the cells 31 of the scheduler 30 have been consulted in turn, the pointer or its equivalent returns to the beginning again.

It will be apparent that when services with different frame formats and lengths are to be scheduled at predetermined constant intervals over the same link, there will be incidences when permits for two or more different services will be scheduled in the same time slot. In this case, the central node 10, or rather
20 its controller, will allocate the time slot to one of the services and cause the frames for the remaining services to be delayed by one or more time slots. The allocation of time slots for conflicting permit schedules may be decided on the basis of prioritising information relating to the services, whereby one service will consistently be handled before the others. Alternatively, the central node
25 10 may treat all services with equal priority by giving one service priority in the first incidence of conflict and another service priority in the next incidence. Alternatively, the scheduler 30 could be programmed such that only one permit is allocated to each cell, however this would necessarily introduce a firm priority for one or other of the services.

5 The scheduler 30 described with reference to Fig. 2 can deal efficiently with the access of circuit emulated delay sensitive traffic over a shared link of a packet switched network when the number of services with different frame lengths is not too high. However, when several different services, many with different frame lengths are to be scheduled on a shared medium, the length of the scheduler will become prohibitively large. Furthermore the scheduler 30 of Fig. 2 does not lend itself easily to expansion. It is complex to incorporate additional services, since new permit schedules will effect both the total length and the content of the scheduler 30.

10 This problem is overcome by the scheduler 40 illustrated in Fig. 3 according to a second embodiment of the invention.

15 In place of the single circular memory of the first embodiment of the invention, this scheduler 40 includes several short schedule memories 41 arranged in a queue. Each schedule memory 41 contains information for services having the same frame format and length. In Fig. 3 the topmost schedule memory 41 contains scheduling information for the service A, the next schedule memory 41 in descending order contains information relating to service B, and the final schedule memory 41 contains permit scheduling information for service C. The maximum length required for this scheduler is thus $a+b+c$ cells 42. Furthermore, since several services could be scheduled using each single queue, provided they have the same frame length the total length of the scheduler will not change. The cells 42 of each schedule memory 41 represent coincident time slots. As indicated by the arrows in Fig. 3, the central node 10, or its controller, thus consults the cells 42 of each of the schedule memories 41 representing the current time slot in turn, and determines which, if any, of the services is to be sent a permit for the specified

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time slot.

As for the scheduler structure 30 of the first embodiment, some conflict is inevitable between the scheduled permits. Accordingly, when permit
5 allocation markers are present in more than one cell 42 representing the current time slot, the central node controller will decide which service will be handled preferentially. This is done on the basis of prioritising information obtained previously from the corresponding multipoint nodes 20.

Alternatively, when no service is defined as having priority, or all services
10 stipulate similar delays and delay variations to be respected, priority will be allocated by the central node controller to each service in turn as each conflict arises. Also, as for the previous embodiment, the schedule memories 40 may be programmed with a specified priority hierarchy causing some services to be consistently delayed in favour of another so that decisions do not need to be
15 made by the central node 10, or its controller.

The scheduling of the transmission of some delay sensitive services over a shared link in a packet switched network in the manner described above can lead to good efficiency and allow the delays and delay variations to be kept
20 small. However, the delays and delay variations can be reduced still further if the scheduling of a service is arranged to correspond with the assembly of cells. Circuit emulated traffic transmitted over a packet switched network may use only part of a cell for traffic information so as to reduce the time needed for assembly of the cell. For example, to reduce the time needed to assemble
25 an ATM cell for a 64kbps telephone service, the ATM cell may contain only 24 bytes of traffic information with the 24 remaining bytes being filled with dummy data. The delay experienced by this service can be substantially reduced, when the scheduling of packet transmission using the scheduler 30, 40 described above is synchronised to coincide with the assembly of the cell

after the arrival of 24 bytes of information. Furthermore, the permits may be scheduled to prevent time slots being wasted on the transmission of idle or dummy cells or cell portions.

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Claims:

1. A method of controlling the transmission of delay sensitive and non-delay sensitive traffic from a plurality of network nodes (20) to a central node (10) over a shared medium (100), wherein time slots for transmission over said shared medium (100) are allocated by said central node (10) to the network nodes (20), characterised by providing a scheduler (30, 40) defining a predetermined scheduling interval for at least one delay sensitive traffic service, said scheduling interval corresponding essentially to the time in terms of time slots for assembling a data packet of the delay sensitive traffic service, prior to allocating a time slot, consulting said scheduler (30, 40) to determine if a scheduling interval has elapsed, and allocating said time slot to a network node (20) carrying a delay sensitive traffic service if said scheduler indicates that the predetermined interval for the delay sensitive traffic service has elapsed.
2. A method as claimed in claim 1, characterised by for each delay sensitive traffic service, adapting the predetermined scheduling interval to the frame length of the traffic service.
3. A method as claimed in claim 1 or 2, characterised by allocating said time slot to a further network node (20) carrying a non-delay sensitive service on the basis of stored information concerning non-delay sensitive traffic awaiting transmission if the scheduler (30, 40) indicates that no scheduling interval has elapsed.
4. A method as claimed in any previous claim, characterised by

determining from said scheduler (30, 40) if the scheduling interval for more than one traffic service has elapsed, and allocating consecutive time slots to the scheduled traffic services with alternating priority.

- 5 5. A method of controlling the transmission of delay sensitive and non-delay sensitive traffic from a plurality of network nodes (20) to a central node (10) over a shared medium (100), wherein traffic is transmitted from said network nodes (20) to said central node (10) in time slots in response to permits issued to said network nodes by said central node characterised by
- 10 providing a scheduler defining a predetermined scheduling interval for at least one delay sensitive traffic service,
- 15 prior to issuing a permit, consulting said scheduler to determine if a predetermined scheduling interval has elapsed,
- 20 if at least one interval has elapsed, issuing a permit enabling transmission during a future time slot to a network node (20) carrying a delay sensitive traffic service, and
- 25 if no interval has elapsed, issuing a permit enabling transmission during said future time slot to a further network node (20) carrying a non-delay sensitive service on the basis of stored information relating to non-delay sensitive traffic awaiting transmission.

6. A packet switched communication network including a central node (10) and a plurality of network nodes (20) connected to said central node (10) by a shared medium (100), wherein data is transmitted from said network nodes (20) to said central node (10) over said shared medium in time slots allocated by said central node, characterised by scheduling means (30, 40) for storing scheduling intervals relating to delay sensitive traffic services carried by at least one network node,

wherein said stored scheduling intervals correspond essentially to the time in terms of time slots for assembling a data packet of the traffic service, and

means in said central node (10) for consulting said scheduling means to determine whether a scheduling interval has elapsed and for allocating a time slot to a network node (20) carrying a delay sensitive traffic service if a scheduling interval relating to said traffic service has elapsed.

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10 7. A network as claimed in claim 6, characterised in that said scheduling means (30, 40) includes a plurality of storage cells (31,42) each cell corresponding to a transmission time slot on said shared medium and being programmable with a marker corresponding to a delay sensitive traffic service.

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8. A network as claimed in claim 7, characterised in that the cells (31, 41) in said scheduling means (30, 40) are sequentially accessible.

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9. A network as claimed in claim 7 or 8, characterised in that the cells of the scheduling means are accessed cyclically.

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10. A network as claimed in any one of claims 7 to 9, characterised in that the number of cells included in said scheduling means (30) is at least equal to the multiple of the scheduling intervals in terms of time slots of the scheduled traffic.

11. A network as claimed in any one of claims 7 to 9, characterised in that said scheduling means (40) comprises several cyclical schedulers (41), wherein each cyclical scheduler (41) is programmed with scheduling

markers relating to traffic services having the same scheduling interval in terms of time slots, the cells of each scheduler corresponding to the same time slot being accessible in turn.

- 5 12. A network as claimed in claim 11, characterised in that each scheduler (41) has a length in cells (42) that is equal to the scheduling interval in terms of time slots of the scheduled traffic services.
- 10 13. A network as claimed in any one of claims 6 to 12, characterised in that said central node (10) comprises queue recording means (11) for storing the packet queue size of non-delay sensitive traffic awaiting transmission at said network nodes (20).
- 15 14. A packet switched communication network including a central node (10) and a plurality of network nodes (20) connected to said central node (10) by a shared medium (100), wherein payload information is transmitted from said network nodes (20) to said central node (10) over said shared medium in time slots allocated by said central node, characterised by
- 20 scheduling means (30, 40) including plurality of storage cells (31,42) each cell corresponding to a transmission time slot on said shared medium and being programmed with markers corresponding to at least one delay sensitive traffic services carried by at least one network node, the number of cells separating a marker for the same traffic service
- 25 corresponding essentially to the time in terms of time slots for assembling a data packet of the traffic service, and means in said central node (10) for sequentially consulting the cells of said scheduling means to determine if a marker designating a delay sensitive traffic service is present and for allocating a time slot to a

network node (20) carrying a delay sensitive traffic service if a marker corresponding to said traffic service is present in a consulted cell.

- 5 15. A packet switched communication network including a central node (10) and a plurality of network nodes (20) connected to said central node (10) by a shared medium (100), wherein payload information is transmitted from said network nodes (20) to said central node (10) over said shared medium in time slots allocated by said central node, characterised by
- 10 scheduling means (30, 40) including a plurality of cyclical storage means (41), wherein each storage means (41) includes a plurality of storage cells (31, 42) each corresponding to a transmission time slot on said shared medium, each storage means further being programmed with scheduling markers relating to delay sensitive traffic services that are carried by at least one network node (20) and have the same frame length in terms of time slots, and means in said central node (10) for consulting the cells (31, 42) of each scheduler corresponding to the same time slot in turn to determine if a marker designating a delay sensitive traffic service is present and for allocating a time slot to a network node (20) carrying a delay sensitive traffic service if a marker corresponding to said traffic service is present in a consulted cell.
- 15 20
- 25 16. A node in a packet switched communication network for controlling the transmission of data from a plurality of network nodes (20) over a shared medium (100) by allocating transmission time slots to said network nodes (20), said node being characterised by means (30, 40) for storing scheduling intervals relating to delay sensitive traffic services carried by at least one network node (20), the scheduling interval for each traffic service corresponding essentially to

the time in terms of time slots for assembling a data packet of the traffic service,

means (11) for recording the number of packets of non-delay sensitive traffic awaiting transmission at said network nodes (20) and

5 a controller for consulting said scheduling means to determine whether a scheduling interval has elapsed and for allocating a time slot to a network node (20) carrying a delay sensitive traffic service if a scheduling interval relating to said traffic service has elapsed or for allocating a time slot to a network node (20) carrying non-delay
10 sensitive traffic selected on the basis of the number of packets awaiting transmission if no scheduling interval has elapsed.

17. A node in a packet switched communication network for controlling the transmission of data from a plurality of network nodes (20) over a shared medium (100) by allocating transmission time slots to said
15 network nodes (20), said node being characterised by scheduling means (30, 40) including plurality of storage cells (31,42) each cell corresponding to a transmission time slot on said shared medium and being programmable with a marker corresponding to a delay sensitive traffic service carried by at least one network node, the
20 number of cells separating a marker for the same traffic service corresponding essentially to the time in terms of time slots for assembling a data packet of the traffic service, and means for sequentially consulting the cells of said scheduling means to
25 determine if a marker designating a delay sensitive traffic service is present and for allocating a time slot to a network node (20) carrying a delay sensitive traffic service if a marker corresponding to said traffic service is present in a consulted cell.

- 5
18. A network as claimed in claim 17, characterised in that the number of cells (31) included in said scheduling means (30) is at least equal to the multiple of the frame lengths in terms of time slots of the scheduled traffic.
- 10
19. A node in a packet switched communication network for controlling the transmission of data from a plurality of network nodes (20) over a shared medium (100) by allocating transmission time slots to said network nodes (20), said node being characterised by scheduling means (30, 40) including a plurality of cyclical storage means (41), wherein each storage means (41) includes a plurality of storage cells (31, 42) each corresponding to a transmission time slot on said shared medium, each storage means further being programmed with scheduling markers relating to delay sensitive traffic services that are carried by at least one network node (20) and have the same frame length in terms of time slots, and means in said central node (10) for consulting the cells (42) of each scheduler corresponding to the same time slot in turn to determine if a marker designating a delay sensitive traffic service is present and for allocating a time slot to a network node (20) carrying a delay sensitive traffic service if a marker corresponding to said traffic service is present in a consulted cell.
- 15
- 20
20. A node as claimed in claim 19, characterised in that the number of cells (41) separating a marker for the same traffic service corresponds essentially to the time in terms of time slots for assembling a data packet of the traffic service at the network node (20).
- 25

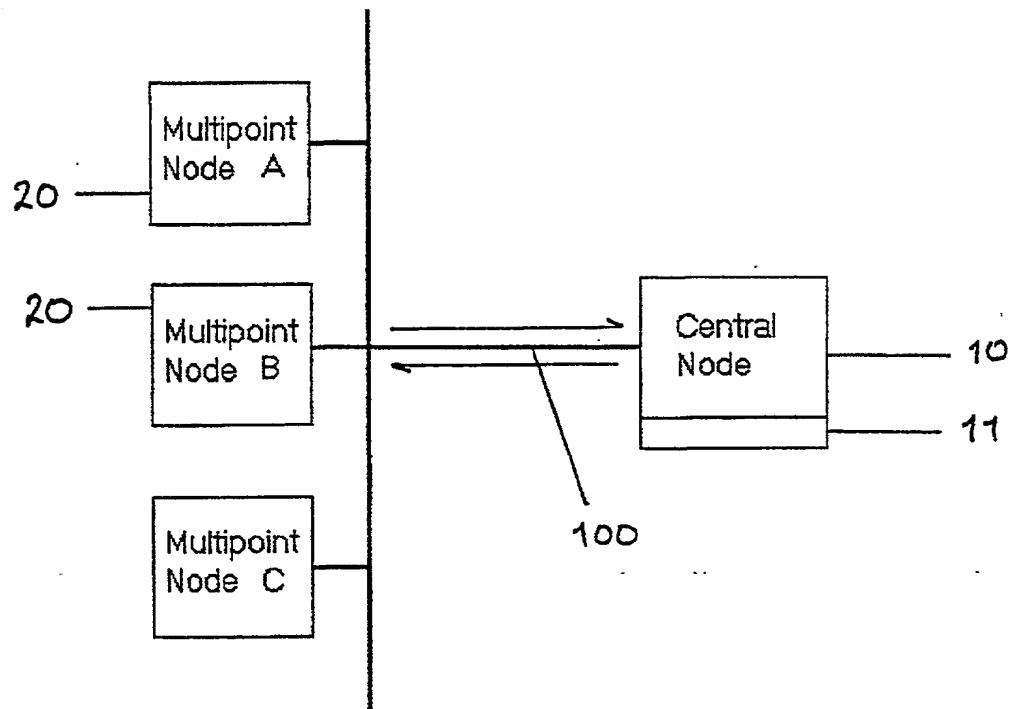


Fig. 1

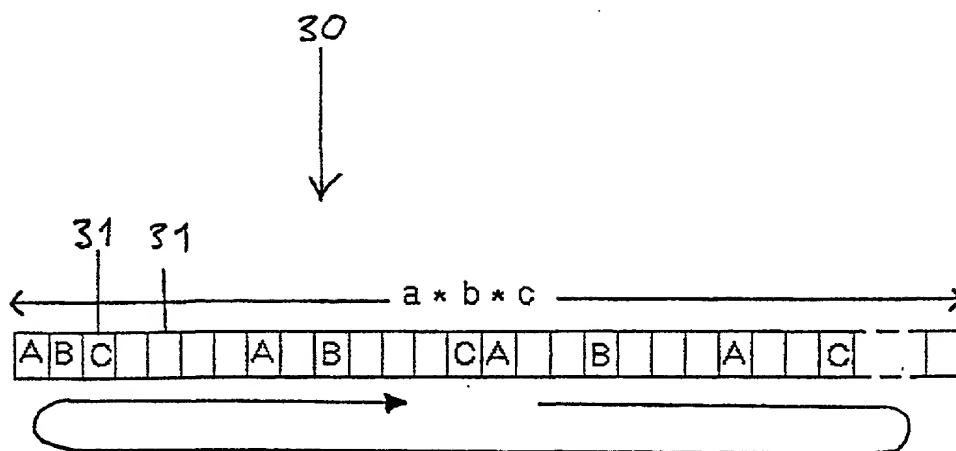


Fig. 2

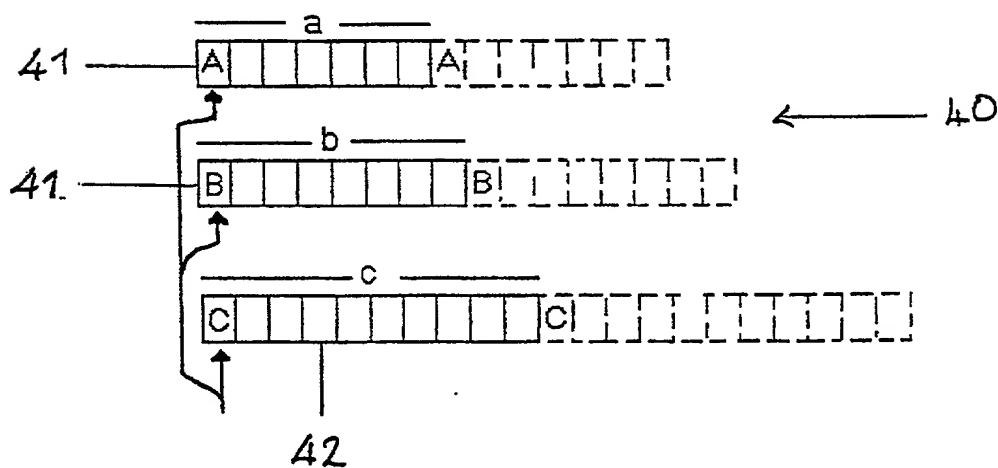


Fig. 3

COMBINED DECLARATION AND POWER OF ATTORNEY

As a below named inventor, I hereby declare that

My residence, post office address and citizenship are as stated below next to my name.

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled:

ACCESS CONTROL MECHANISM FOR PACKET SWITCHED COMMUNICATION NETWORKS

the specification of which: *(check one)*

REGULAR OR DESIGN APPLICATION

☐ is attached hereto.

☒ was filed on June 1, 2001 as application Serial No. and was amended on (if applicable).

PCT FILED APPLICATION ENTERING NATIONAL STAGE

☒ was described and claimed in International application PCT/SE99/02244 filed on 1 December 1999 and as amended on (if any).

I hereby state that I have reviewed and understand the contents of the above-identified specification, including the claims, as amended by any amendment referred to above.

I acknowledge the duty to disclose information which is material to patentability as defined in Title 37, Code of Federal Regulations, §1.56.

PRIORITY CLAIM

I hereby claim foreign priority benefits under 35 USC 119 of any foreign application(s) for patent or inventor's certificate listed below and have also identified below any foreign application for patent or inventor's certificate having a filing date before that of the application on which priority is claimed.

PRIOR FOREIGN APPLICATION(S)

Country	Application Number	Date of Filing (day, month, year)	Priority Claimed
Sweden	9804172-6	1 December 1998	yes

(Complete this part only if this is a continuing application.)

I hereby claim the benefit under 35 USC 120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of 35 USC 112, I acknowledge the duty to disclose information which is material to patentability as defined in Title 37 Code of Federal Regulations §1.56 which became available between the filing date of the prior application and the national or PCT international filing date of this application:

(Application Serial No.)

(Filing Date)

(Status--patented, pending, abandoned)